

Optimal Offering Strategy of a Price-Taking Virtual Power Plant (VPP)

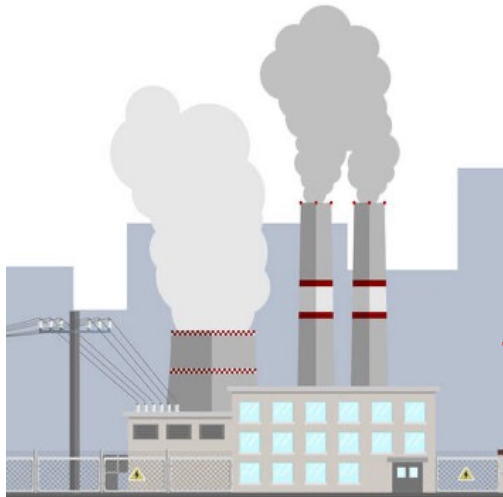
--2023 EEE symposium

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Sep 15, 2023

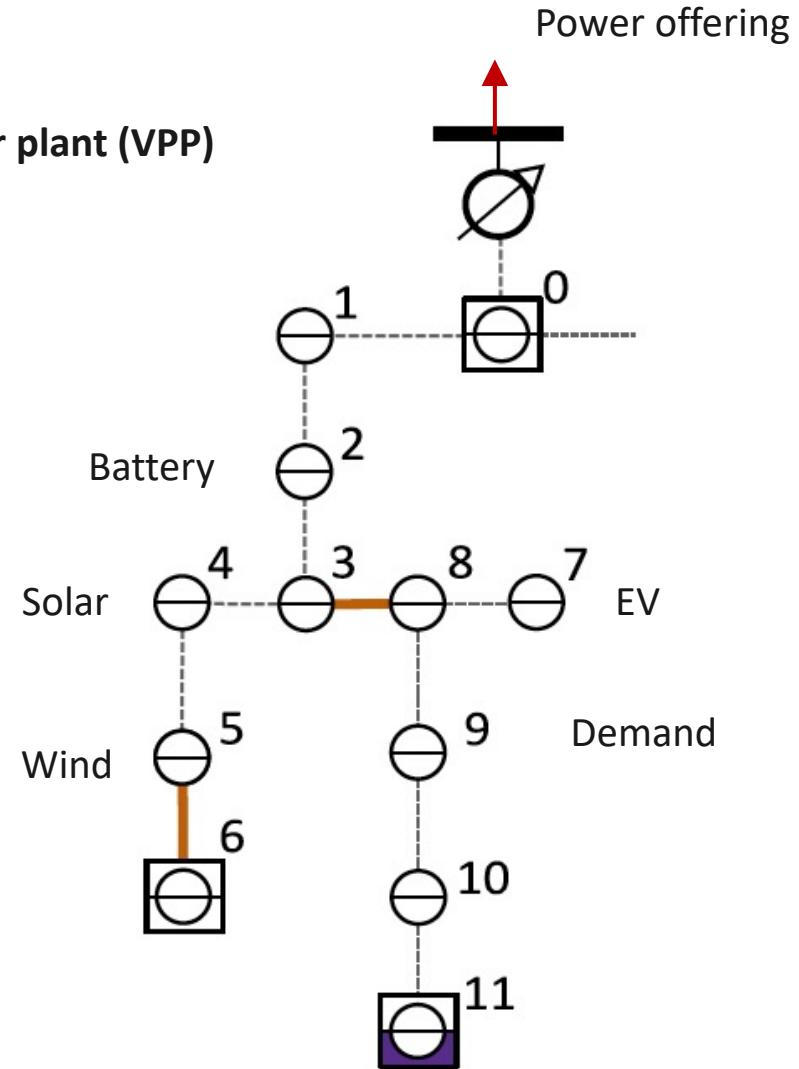
What is Virtual Power Plant

Power plant--Generator



Power offering

Virtual power plant (VPP)



R. Mieth and Y. Dvorkin, "Distribution Electricity Pricing Under Uncertainty,"
in IEEE Transactions on Power Systems, vol. 35, no. 3, pp. 2325-2338, May 2020

What is the significant difference between the VPP
and the traditional generator?

Content

1 Background and contribution

2 Formulation

3 Solution algorithm

4 Results and future works

Background and contribution

Why this topic?

What are the research gaps?

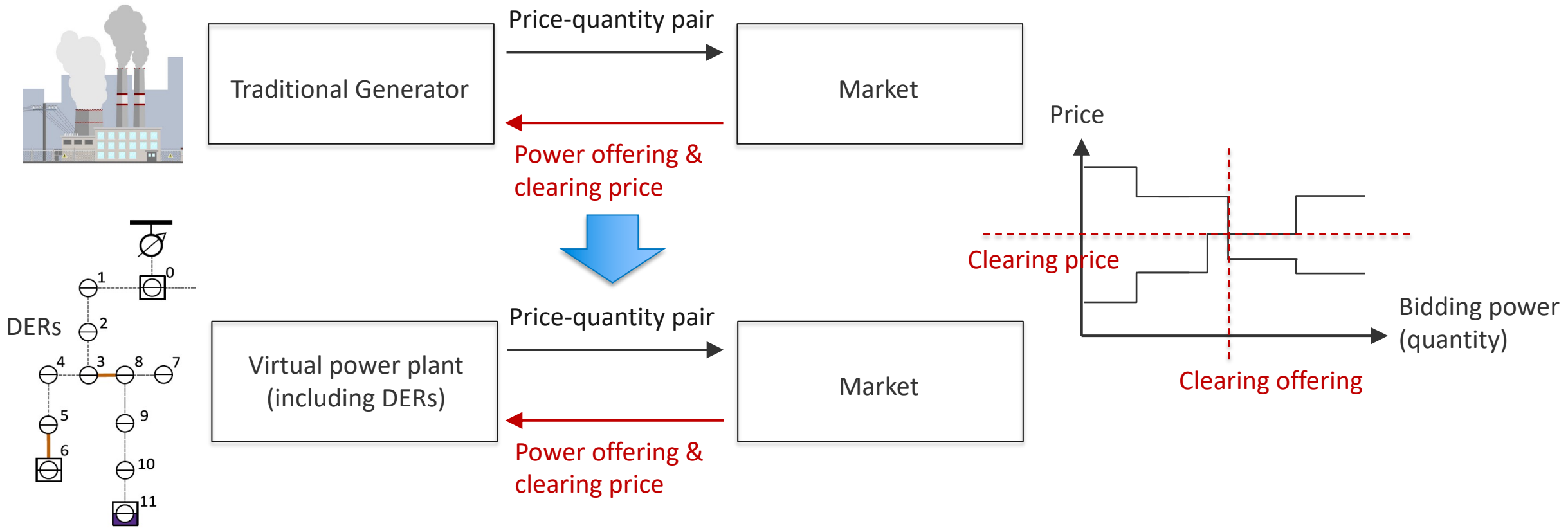
What did we do?

Background

Generator's profit

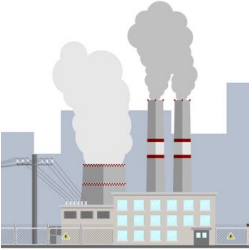
Profit = revenue of selling power + cost of generating power

Day-ahead market



Background

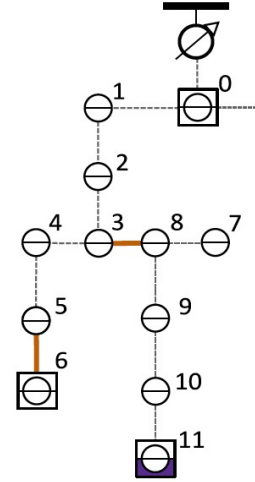
Real-time market



Traditional Generator: Generate required power offering according to the day-ahead market clearing results



Cost is determined by fuel cost



Virtual power plant: Change DERs dispatch after price realization in real-time market

- Providing the same offering power
- Minimize generation cost



Cost is determined by profit maximization considering Real-time market price uncertainty

Gaps & Contribution

Current research

Model price uncertainty under simplified assumptions of stochastic price evolution:

- Uncertainty intervals
- Finite samples of price scenarios



Compromise the optimality of the VPP's power offering

Contribution

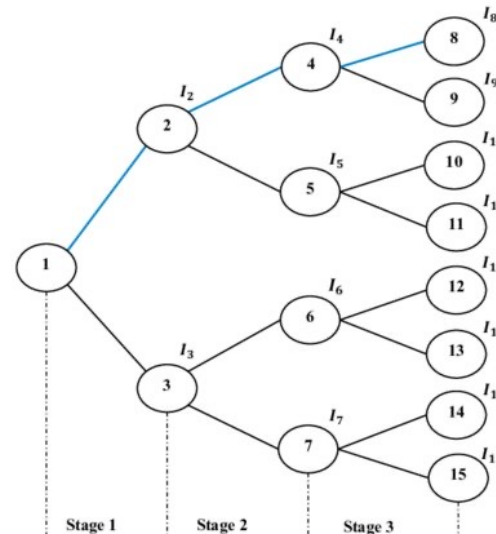
Model: VPP offering strategy is determined by profit maximization

$$\max_x (f(x) - \min_y g_x(y))$$

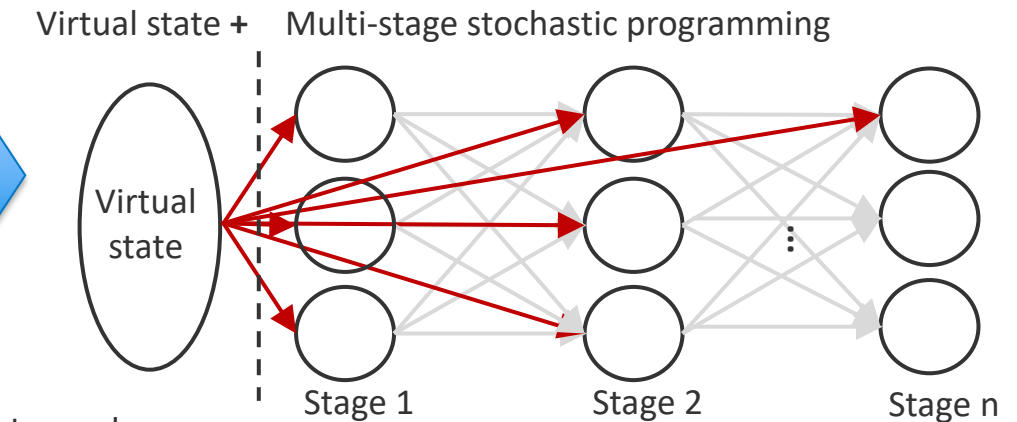
x – power offering
 y – DERs control



Technology challenge: exponential price scenario – curse of dimensionality



Formulation: virtual stage + 24 stages



Legend

→ Power offering

→ State variable

○ A price realization

Formulation

How to formulate the exponential price scenarios?

Formulation

Multi-stage stochastic programming formulation

Simplified

$$\max_x (f(x) - \min_y g_x(y))$$

Revenue Cost

Full

$$\max_{\mathbf{p}} \left\{ \sum_{t=1}^{24} \sum_{s=1}^{|\pi_t|} \mathbb{E}_{\pi_t} [\pi_t^\top \mathbf{p}_t] - \mathbb{E}_{\pi_1} \left[\min_{x_1 \in F_1(p_1)} \{c_1(x_1) - \dots - \mathbb{E}_{\mathbf{p}_{24}|p_{[1,23]}} \left[\min_{x_{24} \in F_{24}(x_{23}, p_{24})} c_{24}(x_{24}) \right] \} \right] \right\}$$

Virtual stage: determine the power offering corresponding to prices, that is to get the price-quantity pair, maximize profit

Subproblem:

$$\max_{p_{1-24,s}} f(p_{1-24,s}) + \sum_{m \in \chi_0} q_{0,m} Q_m(p_{1-24,s})$$

$$f = \sum_{t=1}^{24} \sum_{s=1}^{|\pi_t|} \pi_{ts} p_{ts}$$

Cost-to-go function =
Transition probability *
Optimal value function

Stage 1: determine first time slot DERs dispatch, minimize generation cost

⋮

Stage 24: determine 24 time slots DERs dispatch, minimize generation cost

Subproblem:

$$Q_1(p_{1-24}) = \min_{x_1, y_1} \{C_1(x_1) + \sum_{m \in \chi_1} q_{1,m} Q_m(x_1, p_{1-24})\}$$

$$C = C_{b,s,t} + C_{dg,s,t} + C_{cur,s,t}$$

$$C_{b,s,t} = \lambda^{\text{deg}} (\eta^c p_{s,t}^c + p_{s,t}^d / \eta^d)$$

$$C_{dg,s,t} = a(p_{s,t}^{\text{dg}}) + b$$

$$C_{cur,s,t} = a(p_{s,t}^{\text{cur}}) + b$$

Cost of battery

distributed generator

demand curtailment

S.t. $\sum p_n = L_n$

Power balance

$$p_{\text{dg},\min} \leq p_{s,t}^{\text{dg}} \leq p_{\text{dg},\max}$$

Generation limit

$$p_{\text{dg},\min} - RR \leq p_{s,t}^{\text{dg}} - p_{s,t-1}^{\text{dg}} \leq p_{\text{dg},\max} + RR$$

Ramp rate

$$\text{SOC}_{s,t} - \text{SOC}_{s,t-1} = (\eta^c p_{s,t}^c - p_{s,t}^d / \eta^d)$$

Soc transition

$$0 \leq p_{s,t}^c, p_{s,t}^d \leq p_{b,\max}$$

Charge

$$0 \leq \text{SOC}_{s,t} \leq \text{SOC}_{\max}$$

Soc

$$0 \leq p_{s,t}^{\text{cur}} \leq p_{\text{cur},\max}$$

Curtailment

Network constraints

Solution algorithm

How did we reduce the price scenarios?

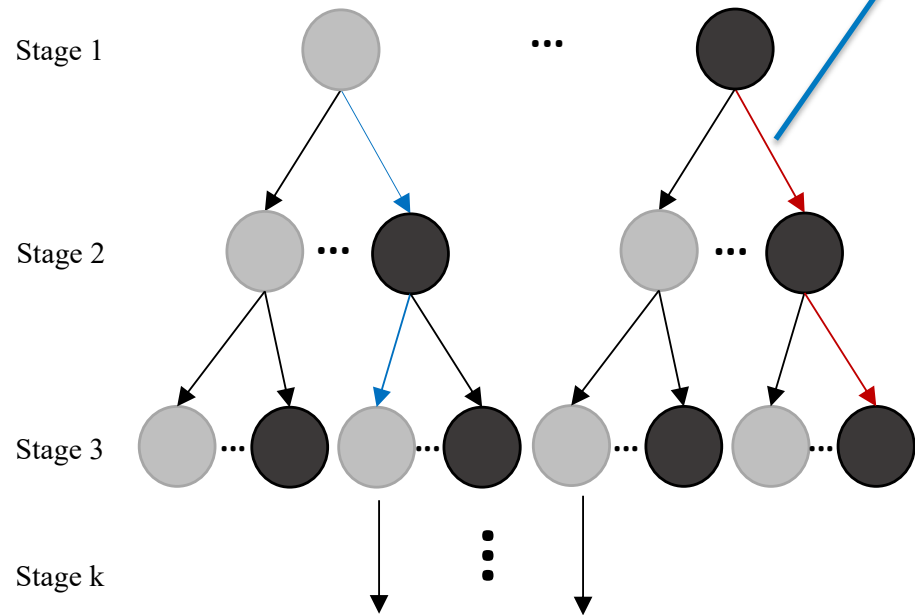
How did we approximate the value function and solve it?

How to deal with the virtual stage

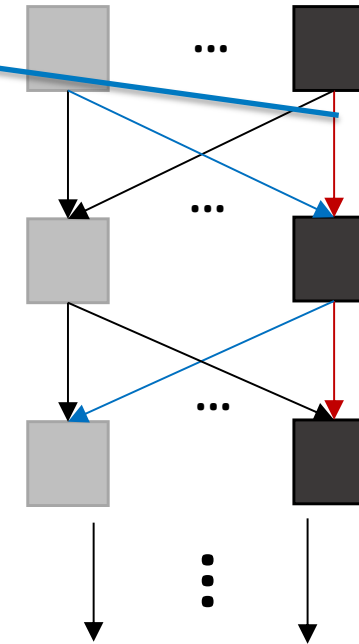
Solution Algorithm

Scenario tree & Policy graph

One line from stages 1-24 indicates a price trajectory, which is exponential.
Corresponding to the transition trajectory in the policy graph



Exponential price scenario



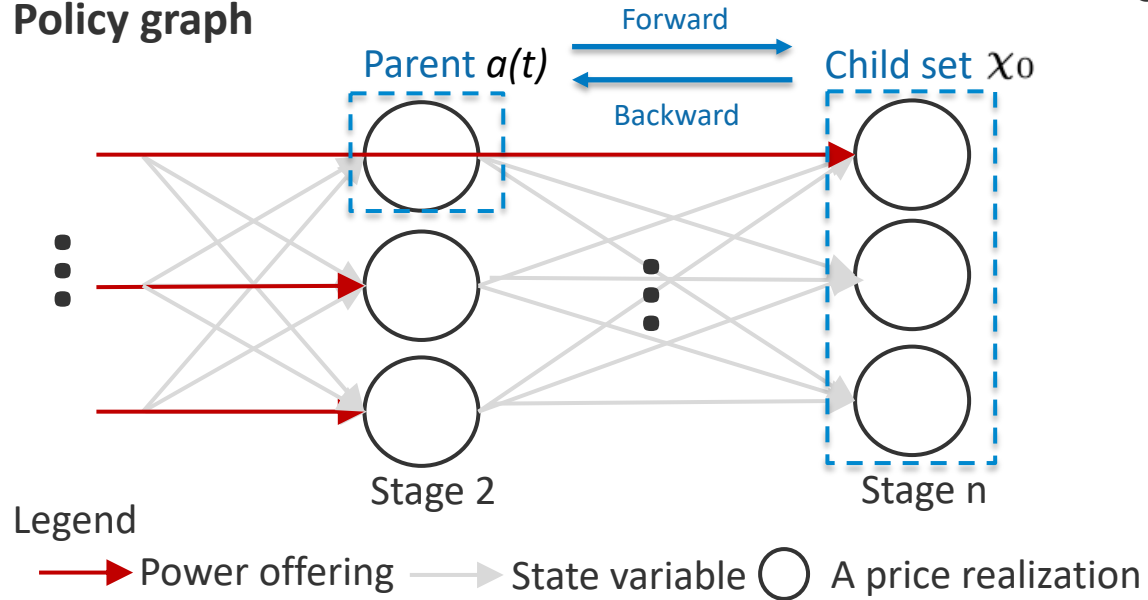
Each node is a price realization, including a decision rule $U=f(x,w)$, give one uncertain realization w , and state variable (time coupled variable) x , and get a different decision U .

Finite policy graph



Solution Algorithm

Policy graph



- The first virtual stage determines the power offering corresponding to prices realization
- Transmit power offering to each corresponding price realization
- When VPP changes DERs dispatch, it should follow the power offering determined in the virtual stage

Solution algorithm – stochastic dual dynamic programming (SDDP)

Each iteration i :

- Forward: Update state variable x_t following $t=1$ to $t=24$. Approximate lower bound of the value function by solving the following optimization model. Passing x_t to child node's (m) problem.
- Backward: Solve suitable relaxation to get cuts following $t=24$ to $t=1$, update the cost-to-go function by adding cuts from all child nodes to parent nodes, change $\psi_{t,i}$ to $\psi_{t,i+1}$, then solve the suitable relaxation of the updated problem again.

Iteration i to $i+1$

Need to know the parent nodes' $a(t)$ state variable

$$\underline{Q}_{t,i}(p_{1-24}, x_{a(t)}, \psi_{t,i}) = \min_{x_t} C_t(x_t, p_{1-24}) + \psi_{t,i}(x_t, p_{1-24})$$

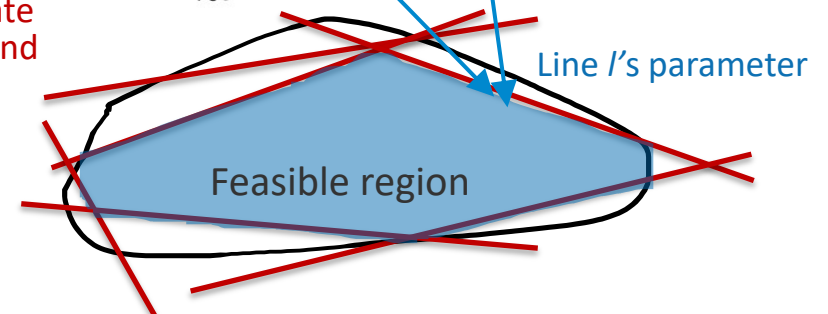
Expected cost-to-go function

$$\text{s.t. } (p_{1-24}, x_t, y_t) \in F_t$$

$$\psi_{t,i}(x_t, p_{1-24}) = \min\{\theta_t : \theta_t \geq LB_t,$$

$$\theta_t \geq \sum_{m \in \chi_0} q_{t,m}(v_{m,l} + \pi_{m,l}^T x_t), \forall l = 1, 2, i-1\}$$

Add cuts update the lower bound



Solution Algorithm

To summarize

- The forward problem is to solve the optimization model to get the optimal solution: state variable, control variable, and cost-to-go function.
- Update statistical upper bound
- The backward is to solve the relaxation problem to get cuts, and add cuts to update (reduce) the feasible region for the forward optimization. (cuts only cut infeasible area)
- Update lower bound

/* (Statistical) upper bound update */

$$\hat{\mu} \leftarrow \frac{1}{M} \sum_{k=1}^M u^k \text{ and } \hat{\sigma}^2 \leftarrow \frac{1}{M-1} \sum_{k=1}^M (u^k - \hat{\mu})^2$$

$$UB \leftarrow \hat{\mu} + z_{\alpha/2} \frac{\hat{\sigma}}{\sqrt{M}}$$

/* Lower bound update */

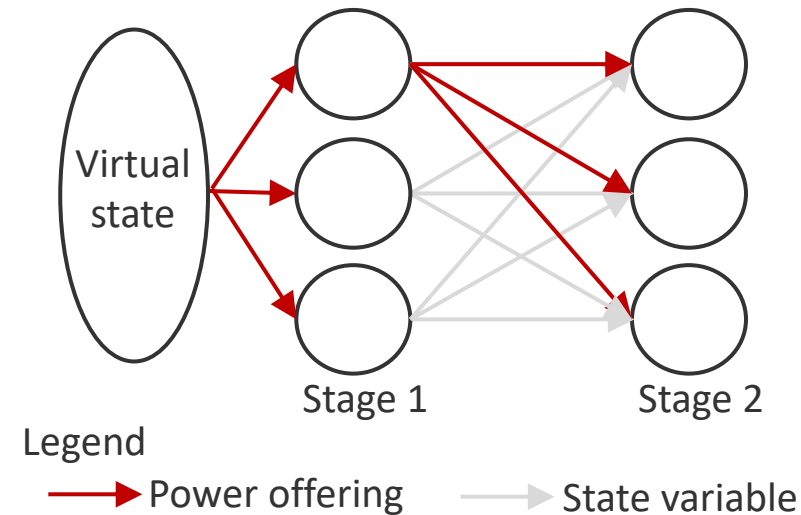
solve $P_1^i(\bar{x}_0, \psi_1^{i+1})$ and set LB to the optimal value

$$i \leftarrow i + 1$$

Optimal value of forward problem at node 1 is the lower bound

Virtual stage – copying variable

Copying power offering variable determined by the virtual stage and transmitted to following stages.



Keep the power offering as a constant and transmit it to next stage step by step

Solution Algorithm

Virtual stage – zeroth order stochastic gradient descent

Algorithm 1 Zeroth-order gradient descent algorithm

- 1: Initialization: perturbation u , number of points for gradient estimation M , initial power offering \mathbf{p} , convergence criteria ϵ , learning rate h
- 2: **repeat** ←
- 3: Call SDDP function ($f(\mathbf{p})$) get optimal DERs dispatch and minimal cost under current offering
- 4: **for** $i=1:M$ **do**
- 5: Randomly generate permutation vector \mathbf{z} , where $\mathbb{E}[\|\mathbf{z}\|] = 1$
- 6: $\mathbf{p}^+ = \mathbf{p} + u * \mathbf{z}$
- 7: $\mathbf{p}^- = \mathbf{p} - u * \mathbf{z}$
- 8: $\nabla f_i = (f(\mathbf{p}^+) - f(\mathbf{p}^-))/2u$ ←
- 9: **end for**
- 10: $\nabla f = -\nabla f_{\text{profit}} + \sum_{i \in M} f(i)/M$
- 11: $\mathbf{p} = \mathbf{p} - h * \nabla f$ ←
- 12: Project \mathbf{p} into box constraint $[p_{\min}, p_{\max}]$ ←
- 13: **until** $\|\nabla f\| \leq \epsilon$
- 14: Get price-power offering pairs considering DERs cost

Create an outer loop using zeroth-order stochastic gradient descent to determine the offering power

$$\max_{\mathbf{p}} \left\{ \sum_{t=1}^{24} \sum_{s=1}^{|\pi_t|} \mathbb{E}_{\pi_t} [\pi_t^\top \mathbf{p}_t] - \mathbb{E}_{\pi_1} \left[\min_{x_1 \in F_1(p_1)} \{c_1(x_1) - \dots - \mathbb{E}_{p_{24}|p_{[1,23]}} \left[\min_{x_{24} \in F_{24}(x_{23}, p_{24})} \text{Cost from SDDP} \} \} \right] \right\} \quad (1)$$

Zeroth-order method calculates gradient

Update decision variables

Projection dealing with box constraint

Results and future works

What is the performance of our method?

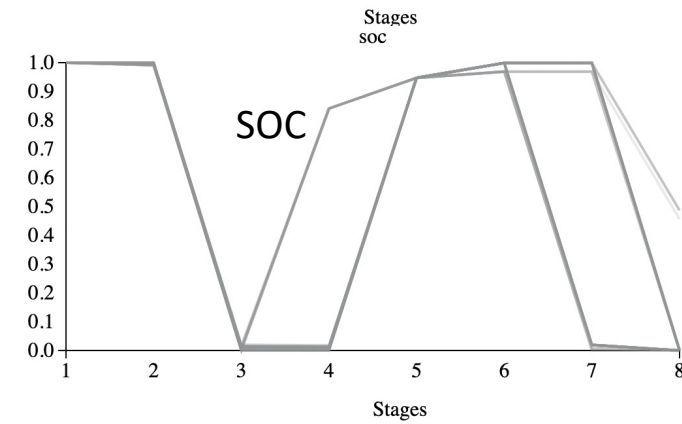
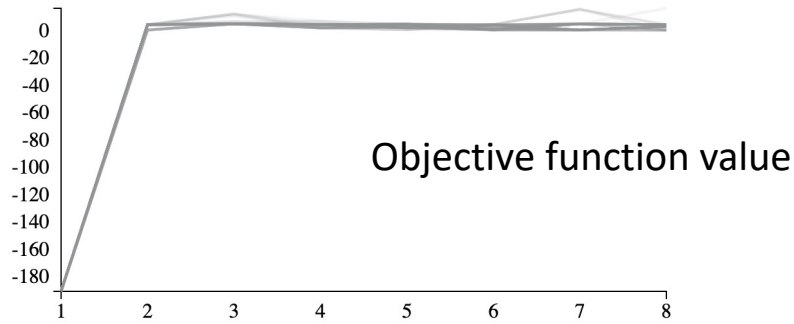
Next step?

Results

Basic setting

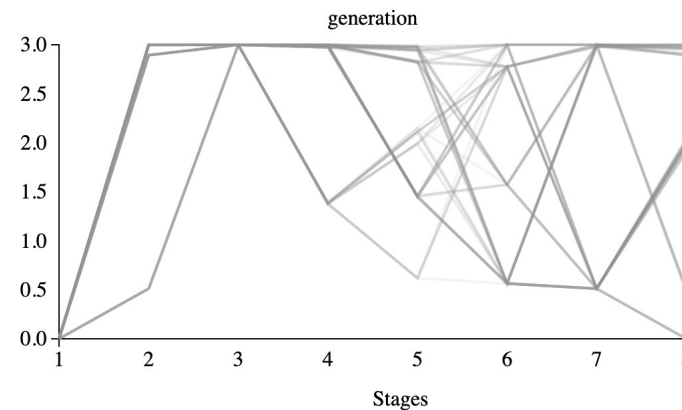
- 1 virtual stage + n stages, each stage includes 5 price scenarios
- 15-bus power network with loads, 6 distributed generators, and 2 batteries
- Using parallel computing and servers to compute

Copying variable - 8 stages (hours)



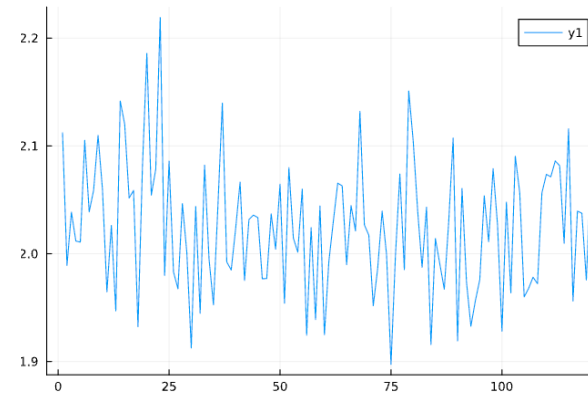
6 minutes running

Distributed generator's generation

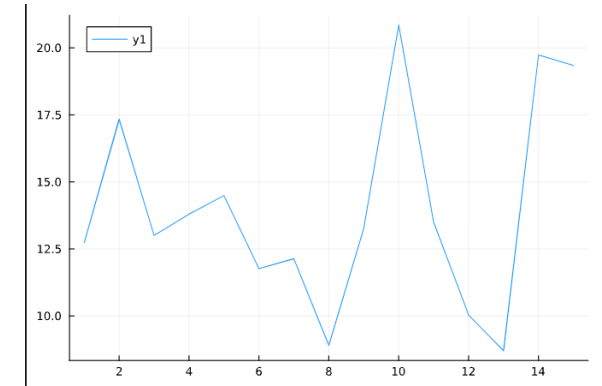


Zeroth-order stochastic gradient descent – 24 hours

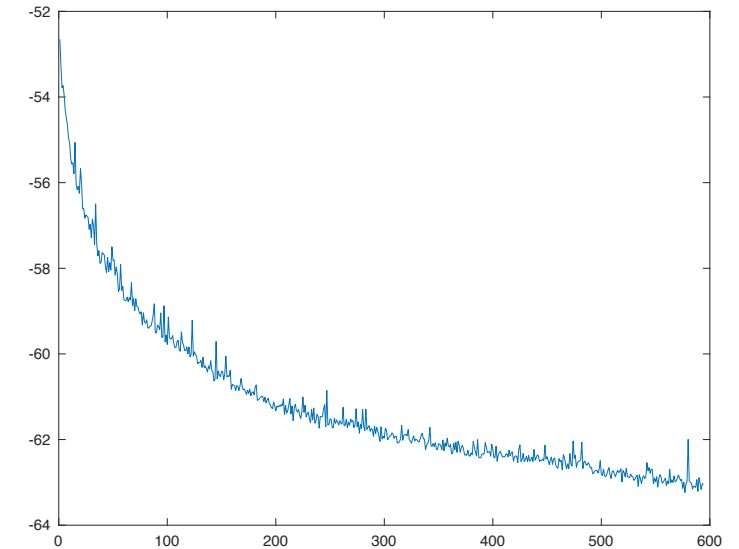
1 hour running



Power Offering



Stochastic gradient recording



Objective function value

Future works

- Update the solution algorithm by adding cuts to the approximate subgradient of the value function, use the proximal bundle method to update the power offering
- Using reinforcement learning to improve the algorithm
- Compare different algorithms' efficiency with regard to this problem
- Including more DERs, and exploring the usage in large power network

Reference

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- Zeroth-order gradient descent. https://scholar.harvard.edu/files/yujietang/files/slides_2019_zero-order_opt_tutorial.pdf

The significant differences between VPP and traditional generators are:

- Stochastic or deterministic generation cost;
- VPP can change DERs dispatch after price realization in the real-time market;
- VPP's offering strategy should consider future price uncertainty
- Multistage stochastic modeling
- Complicated solution algorithm

Thank you

